

TITLE

**A COMMON POWER CONTROL CHANNEL IN A CDMA SYSTEM
AND A SYSTEM AND METHOD FOR USING SUCH A CHANNEL**

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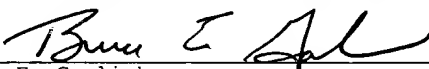
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SPECIFICATION

CROSS-REFERENCES TO RELATED APPLICATIONS

The present application claims priority pursuant to 35 U.S.C. Sec. 120 to U.S. Regular Application Serial No. 09/329,856, filed June 11, 1999, pending, which claims priority under 35 U.S.C. Sec 119(e) to U.S. Provisional Application Serial No. 60/089,154, filed June 12, 1998, and to U.S. Provisional Application Serial No. 60/098,817, filed September 2, 1998, all of which are hereby incorporated by reference in their entirety.

BACKGROUND

1. Field of the Invention

The present invention relates generally to cellular wireless communications and more particularly to the control and operation of common reverse channels in a code division multiple access (CDMA) cellular wireless communication system.

2. Description of the Related Art

Cellular wireless communication systems are generally known to include a plurality of base stations dispersed across a geographic service area. Each of the base stations includes at least one antenna and a base station transceiver system (BTS) and provides wireless service within a respective cell. The BTSs couple to base station

controllers (BSCs) with each BSC serving a plurality of BTSs. Typically, the BSCs also couple to a mobile switching center (MSC) which interfaces to the Public Switched Telephone Network (PSTN) and other MSCs. Together, the BTSs, BSCs and
5 the MSC form a wireless network, which provides wireless coverage to mobile stations (MSs) operating within a respective service area.

Wireless communication systems operate according to various protocol standards. One particular protocol standard
10 in place worldwide is the CDMA protocol standard. CDMA is a direct sequence spread spectrum system in which multiple spread spectrum signals are transmitted and received simultaneously over a common frequency band. In the CDMA system, each mobile station (MS) may be assigned a distinct
15 Walsh code, which identifies the signals, transmitted to and received from the MS.

In an example of operation thereunder, forward link signals from a BTS to a first MS are coded with a first Walsh code and then transmitted where the process of transmission
20 includes pseudo noise (PN) scrambling (spreading). Likewise, forward link signals transmitted from the BTS to the second MS are coded with a second Walsh code and then transmitted, perhaps concurrently with transmissions from the BTS to the first MS. The first MS's receiver receives at its antenna
25 all of the energy transmitted by the BTS. However, because Walsh code channels are orthogonal, after correlating the

received signal with the first Walsh code, the despreader outputs all the energy intended for the first MS but none of or only a small fraction (i.e., due to orthogonality loss) of the energy intended for the second, third, etc., MS.

5 Likewise, the second MS correlates the received forward link signal with the second Walsh code to receive its intended forward link energy. Each of the MSs then operates upon the despread signal energy to extract data intended for the respective MS. The number of users accommodated on the
10 forward link is limited by intra-cell interference due to orthogonality loss, inter-cell interference and other interference such as that due to thermal noise.

In some specialized applications (e.g., fixed access) it may be considered to design a system such that orthogonal
15 codes (e.g., Walsh codes) separate the signals of the reverse link users. However, for all CDMA systems currently deployed and reported in the literature, the reverse link is strictly interference limited, that is, one user's reverse link energy at the base station receiver acts as interference to other
20 users signals. Thus, in the typical case, a plurality of MSs transmit to the BTS simultaneously on the reverse link with each reverse link transmission spread by a unique PN code or PN code shift. A receiver of the BTS receives the composite reverse link signal and despreads the reverse link
25 transmissions with expected PN code to extract signals received from the first, second, third, etc., MS. The BTS

then operates upon each despread signal to extract data sent by the MSs.

Wireless communication systems were originally designed and constructed to service voice communications. However, as
5 packet data communications have increased in popularity, wireless communication systems have been called upon to service not only voice communications but packet data communications as well.

Examples of packet data communications supported by
10 wireless communication systems include Internet sessions, electronic mail transfer, electronic file transfers, and short message services, among other services. Human users of the wireless communication system typically initiate these packet data communications. However, electronic devices may
15 also access the wireless communication system, such devices including vending machines, credit card machines, ATMs and other computer controlled electronic devices. The wireless communication system provides a convenient and cost-effect method for sending packet data transmissions when a landline
20 is not available.

Packet data communications place demands upon wireless communication systems that differ from those placed by voice communications. While voice communications require a substantially constant bandwidth, packet data transmissions
25 are "bursty", with high bandwidth required during some time periods and little or no bandwidth required during other time

periods. In servicing either voice or packet data communications, limited channel resources must be assigned. The setup and initialization of a traffic channel in a CDMA system (for either packet data communications or voice services) requires between 20 and 30 frames. In a typical packet data session, once the traffic channel is setup, packet data transmissions between a MS and a BS will occur for a short period of time to complete the setup of the packet data session at all protocol levels. After the initial setup, the packet data transmissions are typically bursty, with periods of inactivity intermixed with bursts of data. Thus, the traffic channel may become idle for periods of time.

According to traffic channel management operations, when a traffic channel is idle for a period of time, e.g., one minute, the traffic channel is released. Thus, during a packet data session, when the traffic channel is inactive for the period of time, it is released. The release of the traffic channel does not, however, release the logical connection established for the MS via higher protocol layers (i.e., above the physical layer). For example, although the physical connection via the traffic channel may be terminated during periods of inactivity, the IP address of the mobile node, call control, and service information is maintained by the network. When the traffic channel is required again, it is again setup. This operation, in combination with the

bursty nature of packet data communications requires the frequent allocation and release of traffic channels. Because the setup of a traffic channel consumes significant overhead (e.g., 20 to 30 frames of setup information on the traffic channel before transmission of user data), continued allocation and deallocation of traffic channels in servicing packet data communications is undesirable.

Many packet data communications transfer very little data after the initial setup of a logical link between a MS and a remote computer. For example, credit card verification services require the transmission of a relatively small number of frames of data. This amount of data does not justify the reallocation of a traffic channel since more frames would be consumed in setting up the traffic channel than would be used in transmitting the packet data. Thus, subsequent traffic channel allocation after the initial packet data communication setup of a logical link is not justified for this reason as well.

Responsive to these concerns, reverse link channels have been standardized for the transmission of reverse link packet data transmissions (and messages) from the MS to the BS. An example of an interference-based reverse link channel in a CDMA system is the reverse link access channel (R-ACH); a common channel that is defined within the IS-95 standard. The R-ACH employs the well-known Slotted Aloha protocol, where for IS-95 each slot is comprised of a preamble followed by a

message capsule. The preamble is typically 3 to 4 frames and the message capsule is typically 4 to 10 frames. In transmitting packet data on the R-ACH, a MS simply initiates transmission, attempting to successfully complete the transmission to the base station without a collision (i.e., another mobile station trying to send a message during the same slot) from other MSs. Thus, some probability exists that the transmission will be received by the BS without contention from other MSs. The probability of successful transmission to the BS on the R-ACH decreases as the loading increases or with services that require more frequent access via the R-ACH such as telemetry, packet data, or credit card applications. Thus, as usage of the CDMA cellular system increases with new packet data applications the ability to effectively use the R-ACH decreases.

Another problem faced in using the R-ACH relates to the power control of transmissions on the R-ACH. The MS uses the received power of forward link transmissions to estimate the transmission power to use for reverse link transmissions (open loop power control). Based upon its estimate of the transmission power, the MS sends a transmission to the BS on the R-ACH. If the BS does not acknowledge receipt of the transmission, the MS increases its transmission power and retransmits. This process is repeated until the BS acknowledges the transmission or a maximum number of tries have attempted, in which case, the MS ceases its

transmissions. Thus, fast closed loop power control does not exist for the R-ACH and the R-ACH channel requires a higher power for effective operation compared to channels which benefit from fast closed loop power control.

5 Thus, there exists a need in the art for a CDMA system, which has a viable mechanism for satisfying bursty packet data transmissions. Further such mechanism should satisfy the data transmission requirements of a MS without intervention of the BS. According to this mechanism, power control and collision avoidance should be managed on a
10 reverse link channel for support of heavy cell loading and bursty packet communications on the reverse link.

SUMMARY OF THE INVENTION

15 In order to overcome the shortcomings among other shortcomings of prior CDMA systems, a communication system constructed according to the present invention supports designated mode (i.e., or "scheduled mode") data bursts on a reverse common channel from a subscriber unit to a base
20 station. This reverse link common channel supports contention-based and contention-less (i.e., during scheduled or designated mode) transmissions from the subscriber unit to the base station. As contrasted to the prior systems, the subscriber unit may reserve the designated mode data burst
25 via a reservation request message. Further, according to the present invention, reverse link power control during burst mode transmissions is supported to increase throughput of the

reverse link common channel(s).

When transmissions are scheduled and serviced in the designated mode data burst, collisions are avoided. This increases throughput and efficiency of the reverse link common channels. Thus, the reverse link common channel is more efficiently operated. Further, because setting up the designated mode data bursts on the reverse link common channel require little overhead as compared to the setup of a traffic channel the CDMA system is operated at a greater efficiency.

Operation of the subscriber unit and the base station may be substantially compliant with CDMA 2000. In such case, the reverse link common channels may comprise the Reverse Common Control Channel and Reverse Access Channel. Under operation that is substantially compliant with the IS-95 operating standards, the reverse link common channel may comprise the Access Channel.

Designated mode data bursts on the reverse link common channel may consume a single slot or multiple slots. The number of slots consumed in the designated mode data bursts depends upon the volume of data the subscriber unit has to transmit to the base station. Thus, in a designated mode reservation request message (transmitted on the reverse link common channel in a contention-based mode), the subscriber unit may state the amount of data it desires to transmit in the designated mode data burst. Based upon this stated

amount, the base station may reserve multiple slots for the designated mode data burst. Then, the subscriber unit will transmit data during the slots.

The base station may initiate a designated burst mode transmission from the subscriber unit without a reservation request from the subscriber unit. Such operation may be performed, for example when a logical data session was previously set-up and a response from the subscriber unit is expected after the expiration of a particular time period.

Upon the expiration of the time period, the base station will send a designated mode assignment message to the subscriber unit. The subscriber unit will then respond with a designated mode data burst. Similarly, such operation may be performed, for example, after the occurrence of some event in the system.

A common power control channel according to the present invention is a forward link channel from the base station to the mobile stations. The common power control channel supports at least one reverse link common channel and indicates whether the channel has been reserved for a dedicated mode data burst or not. The common power control channel also provides power control bits for mobile stations that are transmitting on the reverse common channel. According to a construction of the common power control channel, a plurality of reverse link common channels are supported with the channel providing reservation indications

for each channel. Further, the common power control channel provides power control bits for each of the plurality of reverse link common channels. Various techniques according to the present invention may be employed to include the various power control and reservation indications within the common power control channel.

Moreover, other aspects of the present invention will become apparent with further reference to the drawings and specification which follow.

BRIEF DESCRIPTION OF THE DRAWINGS

A better understanding of the present invention can be obtained when the following detailed description of the preferred embodiment is considered in conjunction with the following drawings, in which:

FIG. 1 is a system diagram illustrating a portion of a CDMA cellular system in which a plurality of subscriber units contend for reverse link channels, receive reservations for designated transmissions on the reverse link channels, transmit upon the reverse link channels and operate to perform reverse link power control upon the reverse link channels, all according to the present invention;

FIG. 2 is a system diagram illustrating in more detail the structure of a CDMA cellular system constructed according to the present invention and its interconnection with the Public Switched Telephone Network and a packet data network,

in combination, these components servicing data sessions;

FIG. 3 is a system diagram illustrating a subscriber unit and a base station constructed according to the present invention and a plurality of channels upon which the subscriber unit and base station communicate;

FIG. 4 is a block diagram illustrating the structure of a reverse link common channel and a transmit power and inhibit sense multiple access control channel (forward link), both of which are operated according to the present invention;

FIG. 5 is a logic diagram illustrating operation of a subscriber unit in communicating with a base station on a reverse link channel in a common mode;

FIG. 6 is a logic diagram illustrating operation of a subscriber unit in initiating a designated burst mode transmission to a base station and operation during such designated burst mode;

FIG. 7 is a logic diagram illustrating operation of a subscriber unit in transmitting data in a designated burst mode that was initiated by a base station;

FIG. 8 is a logic diagram illustrating operation of a subscriber unit and a base station in performing power control during a reverse link burst mode transmission;

FIG. 9 is a block diagram illustrating the manner in which a base station couples a plurality of power control bits and a plurality of inhibit sense multiple access bits

upon a single forward link channel;

FIG. 10 is a block diagram illustrating three instances of a forward link channel serviced according to the present invention which provides a plurality of power control bits and or a plurality of inhibit sense multiple access bits to a plurality of subscriber units;

FIG. 11 is a block diagram illustrating the manner in which a base station may couple a forward link power control and an inhibit sense multiple access channel to another forward link channel using IQ mapping;

FIG. 12 is a block diagram illustrating generally the components of a base station and of a subscriber unit, each of which operates according to the present invention.

FIG. 13 is a block diagram illustrating generally one methodology for positioning power control and inhibit bits on a common power control signal according to the present invention; and

FIG. 14 is a block diagram illustrating generally another methodology for positioning power control and inhibit bits on a common power control signal according to the present invention.

DETAILED DESCRIPTION OF THE DRAWINGS

FIG. 1 is a system diagram illustrating a portion of a CDMA cellular system 100 in which a plurality of subscriber units 106-122 contend for reverse link channels, receive

reservations for designated transmissions on the reverse link channels, transmit upon the reverse link channels and operate to perform reverse link power control upon the reverse link channels, all according to the present invention. The CDMA cellular system infrastructure shown includes a base station 102 and a network infrastructure 104. These components will be discussed in more detail with reference to FIG. 2 and additional network infrastructure 104 components will be introduced. However, these components are generally known and will be described only as they relate to the teachings of the present invention. The CDMA cellular system 100 may operate according to any various industry standard protocol (or proprietary protocol) that has been modified in accordance with the teachings of the present invention.

The base station 102 provides wireless service within a corresponding geographic area (e.g., cell or sector(s)). The subscriber units 106-122 establish respective wireless links with the base station 102 and, once established, pass voice communications and data communications across the wireless links. As will be discussed further with reference to FIG. 3, the base station 102 and the subscriber units 106-122 support various forward link and reverse link channels. While some of the channels are tailored for substantially constant bandwidth voice communications, other of the channels are tailored to "bursty" data communications.

Thus, some of the subscriber units (e.g., voice

terminals 118, 120 and 122) require primarily voice communication services. Alternatively, other of the subscriber units (e.g., desktop computer 106, laptop computer 108, wearable computer 110, data terminal 112, vending machine 114 and credit card terminal 116) require primarily data communication services.

In servicing voice communications, the base station 102 and a subscriber unit (e.g., voice terminal 118) interact to setup a traffic channel (TCH). The TCH includes a forward link (F-TCH) from the base station 102 to the voice terminal 118 and a reverse link (R-TCH) from the voice terminal 118 to the base station 102. Communications on the F-TCH are encoded with a long Walsh code assigned to the voice terminal 118 and spread with a PN sequence. This encoding and spreading allows the voice terminal 118 to decode the F-TCH communications. R-TCH transmissions may also be encoded with a long Walsh code unique to the voice terminal 118. However, R-TCH transmissions are not always encoded in such a fashion.

In servicing data communications, the base station 102 may initially setup and service a TCH so that a logical connection may be established between a subscriber unit, e.g., data terminal 112, and another digital computer. After a period of inactivity on the TCH, however, the TCH may then be released. Subsequent bursty data communications are then serviced on a reverse common channel such as a Reverse Access Channel (R-ACH) or a Reverse Common Control Channel (R-CCCH).

The R-ACH and the R-CCH are channels for which the subscriber units contend. There may exist multiple R-ACH and R-CCCH channels with each distinguished by a unique PN spreading code. In combination, these channels are referred to as
5 Reverse Common Channels (R-CCHs). The R-CCHs may be used to transmit messages and/or data from a subscriber unit to the base station 102. Of course, the definition of the R-CCH includes channels other than that R-CCCH and the R-ACH that operate according to the present invention.

10 Slots of the R-CCH include a preamble, a message capsule and an idle period. The size of the preamble and message capsule vary for differing channel structures and implementations. However, in one embodiment, the preamble occupies 3 to 4 frames while the message capsule occupies 4
15 to 10 frames. In initiating contention-based transmissions on the R-CCH, a subscriber unit synchronizes to a pilot signal (broadcast from the base station) and then transmits upon the R-CCH. The transmission is successful if the base station 102 is able to decode the transmission. Upon a
20 successful transmission, the base station 102 sends a confirmation of receipt to the transmitting subscriber unit, e.g., data terminal 112, on a Forward Common Control Channel (F-CCCH). Transmissions may be unsuccessful if collisions occur or if the base station 102 does not receive the
25 transmission at a sufficient power level.

According to the present invention, a subscriber unit,

5 e.g., data terminal 112, may reserve a R-CCH for designated mode data bursts. When transmitting on a given R-CCH in this manner (i.e., by reservation and/or designation), contention is avoided for that specific R-CCH channel. In reserving an R-CCH, the data terminal 112 transmits a designated mode reservation request message to the base station 102 on the R-CCH. If the reservation request message is successfully received by the base station 102, the base station may assign one or more slots to the data terminal 112 for a designated mode data burst. In such case, the base station 102 sends a designated mode assignment message to the data terminal 112. The data terminal 112 then waits for the assigned time slot and begins transmission of the data burst on the R-CCH. The data terminal 112 continue its data burst transmission on the R-CCH during the designed mode assignment period. Once the designated mode assignment period expires, common transmission operation continues. However, the data terminal 112 may make further requests for designated mode assignments.

20 The base station 102 may also initiate designated mode operations. In an example of such an operation, the base station 102 desires particular information from credit card machine 116. Previously, the credit card machine 116 established a data link on the TCH with a remote computer during which time validation was performed. After a period of inactivity, the TCH was released. However, the credit

card machine 116 had additional data to collect and transmit. Because the base station 102 is aware of this additional data transmission requirement, it will subsequently receive this data in a designated mode data burst. When the time comes
5 for the data to be received, the base station 102 sends a designated mode assignment message on the F-CCCH to the credit card machine 116. The credit card machine 116, based upon the message, responds in a designated mode data burst with the data it has collect. The data is then relayed to
10 the remote computer.

According to another aspect of the present, closed loop power control is performed during transmissions from a subscriber unit to the base station 102 on the R-CCH. Such power control may be performed not only during designated
15 mode data bursts on the R-CCH but during common reverse mode transmissions on the R-CCH as well. Closed loop power control on the R-CCH increases the success rate for R-CCH transmissions by causing the transmissions to arrive at the base station 102 at a desired signal-to-noise ratio (SNR).

20 According to still another aspect of the present invention, a channel is defined that carries power control bits for at least one R-CCH. This channel may further carry information regarding designated mode data burst assignments on the R-CCH. The structure of this channel will be further
25 described with reference to FIG. 4.

FIG. 2 is a system diagram illustrating in more detail

the structure of a CDMA cellular system constructed according to the present invention and its interconnection with the Public Switched Telephone Network (PSTN) 206 and a packet data network 210. In combination, the components illustrated

5 service data sessions. In a typical CDMA system (or other cellular system), the base station couples to a Base Station Controller (BSC) 202. Other base stations (not shown) also couple to the BSC 202. The BSC 202 (and other BSCs not shown in FIG. 2) couple to a Mobile Switching Center (MSC) 204.

10 The MSC 204, in turn, couples to the PSTN 206 and to a packet data network 210 (e.g., the Internet) via an Interworking Function (IWF) 208.

Communications routed across the PSTN 206 are typically voice communications (i.e., circuit switched) while

15 communications routed across the packet data network 210 are typically data communications (i.e., packet switched). However, the PSTN 206 may also service data communications (e.g., circuit data such as video) while the packet data network 210 may also service voice communications (e.g.,

20 Voice over Internet Protocol communications). Data communications are serviced between the coupled wireless subscriber units, 108, 114, 116 and 118 and various digital devices coupled to the PSTN 206 and the packet data network 210. These digital devices include, for example, a dial-up

25 client computer 216 coupled to the PSTN 206, a web server 212 coupled to the packet data network 212 and a Service Control

Point (SCP) 214 coupled to the packet data network 210. As is generally known, a wide variety of services may be provided to the wireless subscriber units 108, 114, 116 and 118 via connections with these digital devices. These
5 services will not be described herein except to expand upon the teachings of the present invention.

As is generally known, data communications are serviced between digital computers as contrasted to voice communications, which are serviced between human users.
10 These digital computers are operated by humans in some implementations but are stand-alone devices in other implementations. For example, a human user typically operates a laptop computer 108 and provides significant intervention during some tasks like web browsing but provides
15 little intervention during other tasks like file transfers. As another example, a credit card machine 116 requires little human interaction in performing its functions. Further, a vending machine 114, in reporting its status to a remote computer, e.g., dial up client 216, requires no human
20 intervention once it is programmed.

When human intervention/interaction is required, the human user expects a particular response period. Thus, communications of this type are generally serviced upon request. In the case of a laptop computer 108 user
25 establishing a web session with web server 212, the human user requires at least a minimum response rate. Thus, data

communication services between the base station 102 and the laptop computer 108 must be provided at a minimum data rate. In such case, to provide the minimum data rate, a TCH is setup between the laptop computer 108 and the base station and the data communications are serviced on the TCH. However, when the data communications between the laptop computer and the web server 212 become idle for a period of time, the TCH may be released. Then, when the data communication requirements continue, the TCH is reassigned to service the data communications. Thus, the TCH is setup, released and setup again, as data communications require.

Unfortunately, setup of the traffic channel consumes upward of 30 frames of communications between the base station 102 and the laptop computer 108. Thus, according to the present invention, the R-CCH is used for reverse link data communications. These data communications may be transmitted from the laptop computer 108 in a common mode, in a designated mode data burst initiated by the base station 102 or in a designated mode data burst initiated by the laptop computer 108. Thus, when a TCH is not required for the reverse link transmission bandwidth requirements, setup and use of the TCH is avoided in favor of using the R-CCH.

In an example of operation according to the present invention, the credit card machine 116 is used to validate/receive a credit card transaction. Because no land-line is available (or because of the cost associated with the

land-line), the credit card machine 116 employs a wireless link to the base station 102 to form a communication path to the SCP 124. The SCP 124 supports credit card validation/transaction services. The data requirements of the data communications for the transaction may require an initial assignment of a TCH so that the SCP 124 and credit card machine 116 may initialize the transaction. Once these initial data communications are completed, the SCP 214 must validate the credit card and the credit card machine 116 must gather additional information regarding the current transaction. Therefore, because data transmission requirements are reduced, the credit card machine 116 and SCP 214 operate in conjunction with the base station 102 to release the TCH. Then, after a given time period, the SCP 214 requests the additional information from the credit card machine 116.

Because the base station 102 has particular knowledge with respect to the size of this data communication, it requests that the credit card machine 116 respond in a designated mode data burst. The base station 102 therefore sends a designated mode assignment message to the credit card machine 116. In response, the credit card machine 116 transmits its collected data in a designated burst mode on the R-CCH to the base station 102. The base station 102 then relays the data to the SCP 214.

In still another operation according to the present

invention, a vending machine 114 performs a periodic (e.g., daily) notification of its status to a SCP 214. However, this transaction is asynchronous and does not have to be setup between the SCP 214 and the vending machine 114. Thus, at a determined time, the vending machine 114 initiates a communication to the SCP 214 across a wireless link to the base station 102. If the data to sent fits in a single slot, the vending machine 114 performs a common data transmission on the R-CCH. However, if the data to be sent will consume multiple slots, the vending machine 114 sends a designated mode reservation request to the base station 102 on the R-CCH. In response, the base station 102 initiates a designated mode assignment and notifies the vending machine 114 of such assignment. Then, the vending machine transmits its data to the base station 102 in a designated mode data burst during the assigned time period. In either case, the base station 102 receives the data transmission and forwards the data transmission to the SCP 214.

FIG. 3 is a system diagram illustrating a subscriber unit 112 and a base station 102 constructed according to the present invention and a plurality of channels upon which the subscriber unit and base station communicate. A pilot signal is transmitted from the base station 102 and received by all subscriber units (including data terminal 112) operating in a supported geographic area (e.g., cell or sector(s)). Data contained in the pilot signal is used in determining

communication parameters to be used with the base station 102, in synchronizing receipt of transmissions from the base station 102, in synchronizing transmissions to the base station, and in performing open-loop reverse link power control calculations.

The Forward Control Channel (F-CCH) is used by the base station 102 to send control messages and other information to the subscriber units. An example of such a control message is a traffic channel assignment. A F-TCH and a R-TCH are the two components (i.e., forward link and reverse link, respectively) of a TCH, which are setup and serviced to and from the base station 102 and a subscriber unit. The R-ACH and the R-CCCH are examples of R-CCHs that may operate according to the present invention to service reverse link control messages and data transmissions. The R-ACH and the R-CCCH may both be operated according to the present invention to support designated mode data bursts from the subscriber unit 112 to the base station 102. While only a single R-ACH and a single R-CCCH are illustrated, the base station 102 may support multiple R-ACHs and multiple R-CCCHs.

Further illustrated is a Transmit Power Control and ISMA Control Channel (TPICCH) constructed according to the present invention. The TPICCH provides reverse link power control information to subscriber units transmitting on a R-CCH. Such transmissions could be common or designated mode data burst based. In either case, the reverse link power control

information allows the subscriber unit(s) to perform closed loop power control operations on the R-CCH during transmissions. The TPICCH further provides ISMA information corresponding to the R-CCH channel(s) supported by the base station 102. For designated mode data bursts, this ISMA information assigns a particular R-CCH to a particular subscriber unit for one or more scheduled slots. Upon receipt of this ISMA information, a receiving subscriber unit will determine when it may transmit on the assigned R-CCH and for what duration.

FIG. 4 is a block diagram illustrating the structure of a reverse common channel (R-CCH) 400 and a transmit power control and inhibit sense multiple access control channel (TPICCH), both of which are operated according to the present invention. The R-CCH 400 includes a preamble 402, a message capsule 404 and an idle time 406, these components repeated for each slot in the R-CCH 400. The preamble 402, for example, consumes from between a fraction to 4 frames of the slot and allows the base station 102 to identify the transmission. The message capsule 404, for example, consumes from 1 to some designated maximum number of frames and may include a message and/or data. The idle time 402 consumes the remaining portion of the R-CCH 400 slot. Of course, the number of frames contained in each slot and the number of frames consumed by each of R-CCH component may vary from R-CCH to R-CCH, depending upon the particular embodiment and

may vary without departing from the scope of the present invention.

Transmission of the R-CCH 400 slot is substantially synchronized to the forward link pilot (or other timing signal). The accuracy of the synchronized transmission, however, is limited by multi-path fading, transmission delays and other timing based concerns. Multiple R-CCHs may be transmitted in a common frequency band, with the multiple R-CCHs separated by a PN code, PN code offset, time offset, and/or a Walsh code. In the example shown, the first R-CCH 400 is aligned with a 0-offset (within the accuracy possible considering the aforementioned inaccuracies) while a second R-CCH is aligned with a 1-offset. In this example an offset of 1 is symbolic and may correspond to several PN chips.

The TPICCH 420 includes power control information and ISMA information corresponding to at least one R-CCH supported by the base station 102. Subscriber units transmitting on the R-CCH use the power control information for closed loop power control. Listening subscriber units use the ISMA information to determine whether a corresponding R-CCH has been reserved for a data burst. Thus, the bits transmitted by the base station 102 on the TPICCH 420 (i.e., power control and ISMA bits) direct subscriber units to increase/decrease reverse link transmit power levels during reverse link transmissions. Furthermore, the presence of these bits also indicates that the associated R-CCH is in

use, while the absence of these bits indicates that the associated R-CCH is idle. Therefore, the bits transmitted on the TPICCH serve as both power control and inhibit indicators. That is, the presence of the power control bits indicates that other mobiles should not transmit (i.e., should inhibit) their transmissions on the associated R-CCH.

While FIG. 4 illustrates one stream of TPICCH bits, many streams of power control bits may be transmitted on the TPICCH, each stream offset from each other stream. Furthermore many TPICCHs may be transmitted. Each power control bit stream may be associated with a particular R-CCH with the mapping of power control and ISMA bits on the particular TPICCH bit stream mapped to a specific R-CCH. Thus, there is a one to one correspondence between the TPICCH bits and R-CCH. Such mapping will be discussed in more detail with reference to FIGs. 9, 10 and 11.

During the message capsule, the TPICCH 420 bit stream indicates up/down power control bits to the respective R-CCH users. Mobile stations monitor the TPICCH bits corresponding to a given R-CCH in order to detect an "idle" period following the end of transmission on the R-CCH. The idle period is thus defined as an absence of bits on the TPICCH 420 (corresponding to the specific R-CCH). Once the idle period is detected, the mobile stations may contend for the R-CCH channel on the next slot. Note that the BTS can configure an R-CCH such that is used only for designated

bursts. In such a case the channel will not be accessed in a contention-based fashion and thus all transmissions on the channel are scheduled and may be exercised without an idle time. After the next slot, however, mobiles may not contend for the channel until after the next idle period. If a mobile station detects the presence of ISMA bits 406 corresponding to a given R-CCH offset, then it is inhibited from transmitting during the next slot on the corresponding R-CCH. In short, during the message capsule 404, each component of the bit stream associates with a given access channel offset. The bit stream simultaneously provides power control bits 422A-422H and ISMA bits.

Power control performance is very sensitive to the delay from the time that the base station issues a power control command to the time the power control command is executed by the mobile station. Each power control command is interpreted individually and, if a power control command is received in error (i.e., due to fading and noise in the channel), then the power control is corrected during subsequent power control commands and not the current power control command. Thus, the reverse link power control closed loop is tolerant of a modest power control command error rate (e.g., up to approximately 10%).

In contrast, an error in reception of ISMA bits has a much different affect on performance. For example, false detection of ISMA bits during the idle period 406 will cause

under utilization of the access channel. Missed detection of ISMA bits will result in a collision. Consequently a simple forward error control block code may be employed for protection of the ISMA bits.

5 The error control block code may take the form of a repetition code. In such a case, more than one ISMA bit is transmitted during the idle period. Redundancy in transmission of this type provides some performance improvement due to time diversity. Alternatively, the ISMA
10 bits might consist of one bit (or possibly repeated ISMA busy indications) followed by some parity bits.

For example, in one mode of operation, after a block code is decoded, the presence of the ISMA bits may be estimated by the output of the block decoder. For
15 simplicity, this estimate may be normalized, where a 1 indicates presence of the ISMA bits and a 0 indicates that ISMA bits are not present. The decoder may output hard decisions (i.e., either a 1 or a 0) or a soft value (i.e., some quantized value between 1 and 0). A soft decision
20 output allows at least two thresholds - a Zero Threshold, and One Threshold to be defined. The Zero Threshold is a soft value above zero and below the One Threshold. In a statistical sense, if given that the ISMA bits are not present, and the soft estimate is below the Zero threshold,
25 then the decision is 0. Given that the ISMA bits are present, if the soft estimate is above the One Threshold then

the decision is 1. An uncertainty region effectively lies between the Zero Threshold and One Threshold. Thus if the soft decision lies within the uncertainty region, some resolution of this uncertainty can favor either 1 or 0. When
5 the 1 decision is favored, then the decision favors avoiding collisions at the risk of reduced throughput. Conversely, when the 0 decision is favored, a higher collision rate is favored in an attempt to increase throughput. Of course, other error control operations may also be employed.

10 FIG. 5 is a logic diagram illustrating operation of a subscriber unit in communicating with a base station on a reverse link channel in a contention-based mode. The reverse link channel is an R-CCH that operates according to the present invention. Prior to operating on the R-CCH, the
15 subscriber unit and base station may setup and communicate on a traffic channel (step 502). The operations required to setup and operate a traffic channel are generally known. If the traffic channel is optionally setup, the traffic channel will be released after either a period of inactivity or by
20 instruction (step 504). As has been previously described herein, data communications often are bursty transmissions that must be initially setup and then serviced according to actual data transmission requirements. Operation at differing protocol layers will be unaffected by whether the
25 TCH or R-CCH carries reverse link data communications. Thus, after setup of a traffic channel (step 502), the

communication is logically complete and the traffic channel may be released (step 504) without release of the logical connection. Subsequent communications may therefore be serviced according to the steps 506-514. However, subsequent
5 communications could also be serviced by additional TCH operation.

Based upon its reverse link communication requirements, the subscriber unit may transmit a message (or data) on the R-CCH in a common mode (step 506). Transmissions in this
10 common mode are made on the R-CCH 400 illustrated in FIG. 4. Thus, the R-CCH includes a preamble 402 and a message capsule 404. Further, the contention-based transmission on the R-CCH 400 is synchronized to a pilot signal with an assigned offset (e.g., offset 0, offset 1, etc.). During the transmission on
15 the R-CCH 400, the subscriber unit (and the base station) may perform closed loop power control operations that are further described with reference to FIG. 8. Adjustments to the subscriber unit's transmit power will be made based upon power control bits contained on a forward link transmission
20 (e.g., TPICCH).

After its common transmission on the R-CCH, the subscriber unit waits for a response from the base station (step 508). If the response interval expires (step 510) to indicate that the base station did not receive the
25 transmission, and a time-out has not occurred (step 512), the subscriber unit retransmits the message (or data) to the base

station (step 508). However, if a response time-out has occurred or if the mobile station receives a confirmation from the base station of receipt of the message (or data), operation ends.

5 FIG. 6 is a logic diagram illustrating operation of a subscriber unit in initiating a designated burst mode transmission to a base station and operation during such designated burst mode. As was previously described, a subscriber unit may request a designated burst mode
10 transmission via the R-CCH when the R-CCH is available (i.e., not reserved for another subscriber unit). In such case, the subscriber unit sends an access probe to the base station on the R-CCH requesting a designated burst mode (step 602). In its access probe, the subscriber unit may indicate the amount
15 of data it has to transmit and may request a particular time for transmission. The access probe also includes an identification of the subscriber unit and may reference a logical connection that has been previously established.

Once the request is transmitted, the subscriber unit
20 waits for a response from the base station regarding the request (step 604). Such response sent to the subscriber unit on the F-CCCH or another forward control channel. If a response interval expires without a response to the request (step 606) and a time-out interval has not expired (step
25 608), the subscriber unit sends another access probe requesting a designated burst mode. However, if the base

station responds to the subscriber unit by assigning a designated burst mode to the subscriber unit (step 610), the subscriber unit decodes the response to determine when to initiate its burst mode transmissions on the R-CCH. The subscriber unit then transmits its designated mode data burst during the designated slot (step 612). If the current designated slot was not the last slot designated (step 616), operation returns to step 612 wherein the subscriber unit waits for the next slot to commence and again transmits. However, if the current slot was the last designated slot (step 616) or if a time-out has occurred (step 608), operation ends. As was the case with FIG. 5, the transmission operations in FIG. 6 during the access probe (step 602) as well as during the designated burst mode (step 612) may include closed loop power control operations that will be further described with reference to FIG. 8.

FIG. 7 is a logic diagram illustrating operation of a subscriber unit in transmitting data in a designated burst mode that was initiated by a base station. As was previously described, three modes of reverse link transmissions on the R-CCH are possible. The first mode of reverse link transmission is the contention-based transmission (described with reference to FIG. 5). The second mode is a designated mode transmission that is initiated by a subscriber unit (described with reference to FIG. 6). The third mode is a designated burst mode transmission that is initiated by the

base station. All three modes may be scheduled to occur after a particular time period upon which both the base station and the subscriber unit agree

Still referring to FIG. 7, in an idle state, the subscriber unit listens to forward link transmissions from the base station (step 702). These forward link transmissions may be received on the F-CCCH, the TPICCH or on another forward link channel. When a designated burst mode transmission is scheduled for the subscriber unit (step 704), the subscriber unit determines the R-CCH, the designated slot(s) and the transmission time period(s) for the designated mode transmission based upon the contents of the forward link transmission.

Based upon its determination, the subscriber unit tunes to the allocated channel and waits for the initial allocated slot (step 708). When the initial allocated slot commences (step 710), the subscriber unit begins its designated burst mode transmission (step 712). Upon commencing its transmission, the subscriber unit determines whether the slot upon which the current transmission was made (step 714) was the last designated slot. If not, the subscriber unit transmits during the next allocated slot (step 714). If so, the subscriber unit returns to the idle state awaiting other transmissions on the forward control channel.

FIG. 8 is a logic diagram illustrating operation of a subscriber unit and a base station in performing power

control during a reverse link burst mode transmission. When not transmitting on the R-CCH, power control operations at the mobile station are idle and the subscriber unit waits for the initiation of a reverse link transmission (step 802).

5 When a reverse link transmission is initiated (step 804) but prior to transmission on the R-CCH, the subscriber unit measures the strength of a forward link pilot signal from the base station (step 806). Based upon the measured strength of the forward link pilot signal, the subscriber unit performs
10 open loop power control by setting its reverse link transmission strength based upon the strength of the pilot signal step 808).

Once the open loop transmission strength has been determined, the subscriber unit commences its transmission on
15 the R-CCH (step 810). During the transmission, however, the subscriber unit receives one or more power control bits from the base station on a forward link channel, e.g., F-CCCH, TPICCH or another forward link channel (step 812). Based upon the state of the power control bits, the subscriber unit
20 may adjust the power of its reverse link transmissions (step 814). With the transmission power adjusted, the subscriber unit continues transmitting on the R-CCH (step 810) until the transmission ends (step 816), at which point, power control operations at the mobile also end.

25 FIG. 9 is a block diagram illustrating the manner in which a base station couples a plurality of power control

bits and or a plurality of inhibit sense multiple access bits upon a single forward link channel. As was previously described, it is desirable to use a single forward link channel to transmit power control information and ISMA information for a plurality of R-CCHs. In the structure of FIG. 9, three power control bit and ISMA bit streams are coupled onto a single forward link channel. This forward link channel will support from 0 to 16 of these bit streams at a TPICCH rate of 19.2 Kbps.

Each switch 906, 908 and 910 receives two inputs: (1) a decimated long code with the addition of an offset in time (in units of PN chip time), and (2) a power control and ISMA control bit stream. For example, assume that there are 24 bits in a power control group period (PCG). Then, the long code plus offset is taken modulo 16 (i.e., largest power of 2 equal to or smaller to 24). Each switch is closed during the period of time (i.e., or bits) during each PCG indicated by the decimated long code plus offset. Each power control and ISMA control bit stream is associated with a specific reverse common channel and or corresponding reverse common channel offset. This association can be established through overhead or system configuration messages sent to the subscriber unit, or can even be fixed definitions. Alternately, the association may be sent as part of a message for the case of Designated/Scheduled mode operation.

The basic principle of operation for an individual power

control and ISMA bit stream is similar to the placement of power control bits within a power control group (PCG) for an IS-95 forward traffic channel. However, the illustrated structure differs in that bits in the TPICCH are not punctured onto another channel and in that the TPICCH can support many power control and ISMA bit streams. For a 19.2 KBPS TPICCH rate there are 24 possible bit locations within the power control group. Sixteen of these locations are used for sending power control bits, where the location of each power control and ISMA bit within a PCG is indicated by the decimated long code state plus offset.

As shown, a long code generator 902 output is decimated by decimator 904 to provide an 800 Hz input that is combined with offsets, the combination of these signals indicating for which bits during the PCG that the switch is closed. Note that the offset shown in FIG. 9 for each switch is, in principle, a different and unique offset. These offsets are not necessarily the same offset by which R-CCH channels are offset on the reverse link. The PCB/ISM bits input to each switch also enter at an 800 Hz rate and therefore the corresponding switch simply acts as a gate passing through the PCB/ISMA bit during the desired time. The offset added to the decimated long code is modulo 16 so that if the offset plus decimated long code exceeds 16 it wraps around through the beginning of the PCG. Thus, according to this technique, each power control and ISMA bit stream is placed pseudo

randomly in the PCG and has a deterministic and unique offset with respect to other power control and ISMA bit streams. The top most bit stream produced by switch 906 has an offset of 0 (bits indicated with label "A"), the middle bit stream
5 produced by switch 908 has an offset of 3 (bits indicated with label "B"), and the lower most bit stream has an offset of 6 (bits indicated with label "C").

While the offsets illustrate in FIG. 9 are shown to be fixed, as will be further described with reference to FIGs.
10 10, 14 and 15, these offsets may, in general, vary over time

FIG. 10 is a block diagram illustrating three examples of the construction of a forward link channel serviced according to the present invention which provides a plurality
15 of power control bits and a plurality of inhibit sense multiple access bits to a plurality of subscriber units. For a 19.2 KBPS TPICCH rate there are 24 possible bit locations within the power control group. Sixteen of these locations are used for sending power control bits, where the location
20 of each power control and ISMA bit is indicated by the decimated long code state plus offset.

The offset for the TPICCH shown in FIG. 10 produced by the decimator 904 and long code generator 902 of FIG. 9 is, in principle, a different offset than the access channel
25 offset. In other words, it is a decimated long code offset (forward link) as compared to an access channel offset

(reverse link). The offsets added to the decimated long code for the PCB/ISMA bits are modulo 16 so that if any offset plus the decimated long code exceeds 16 it wraps around through the beginning of the next PCG. Thus, according to
5 this method each power control and ISMA bit stream is placed pseudo randomly in the PCG and has a fixed offset (modulo 16) with respect to other power control and ISMA bit streams.

Example 1002 shows the decimated long code state corresponding to offset 1 so that bit A appears at bit
10 position 1. Further, in example 1002, the bit positions for the bit streams B and C with offsets of 3 and 6, respectively, occur at bit positions 4 and 7, respectively. Example 1004 shows a decimated long code state corresponding to offset 10 so that bit A appears at position 10. Further,
15 in example 1004, bit B with an offset of 3 appears at position 13 while bit C with an offset of 6 appears at position 0. Example 1006 shows a decimated long code state corresponding to offset 5 so that bit A appears at position 5. Further, in example 1006, bit B with an offset of 3
20 appears at position 8 while bit C with an offset of 6 appears at position 11.

The structure of FIG. 9 and the operations of FIG. 10 can accommodate other bit rates with simple modifications. For example, at a 38.4 KBPS rate the TPICCH could support at
25 least 0 to 32 such bit streams. Alternatively, 16 bits could be transmitted such that they utilize the entire PCG (i.e.,

12.8 KBPS rate, increasing the period of each bit), which would effectively increase the reliability of the power control and ISMA bits for the same transmitted amplitude level. Another possibility is to utilize all 24 bit locations in the PCG of the bitstream.

Further to Fig. 9, in general, several methods exist for mapping the ISMA/TPC bits on the TPICCH. For example, Fig 9, maps the ISMA/TPC bits to the TPICCH according to the designated long code and fixed offsets. The following is general categorization of the schemes for mapping ISMA/TPC bits to the TPICCH.

- Method 1: Place the ISMA/TPC bits in fixed positions of the TPICCH.
- Method 2: Use pseudo random hopping to place the ISMA/TPC bits over 2^n positions, as in the IS-95B standard, leaving additional positions in the TPICCH unused (as in Fig. 9).
- Method 3: Use pseudo random hopping to place the ISMA/TPC bits over 2^n positions and use fixed mapping in the remaining positions.
- Method 4: Use a uniform counter to place the ISMA/TPC bits over all available positions.
- Method 5: Use pseudo random hopping to place the ISMA/TPC bits over all available positions.

For a 9.6 KBPS TPICCH rate, 12 possible bit locations within a PCG are available for sending TPC and ISMA bits.

According to the previously described general methodologies, the following specific implementation may be used:

- Method 1: Twelve (12) CCHs are supported and assigned fixed bit positions.
- 5 • Method 2: Eight (8) CCHs are supported and assigned pseudo random positions using 3 bits derived from the long code mask (similar to Fig. 9).
- Method 3: Twelve (12) CCHs are supported. The bit position corresponding to 8 CCHs are pseudo randomly changing using 3 bits derived from the long code, while the remaining 4 channels are assigned 4 fixed positions.
- 10 • Method 4: Twelve (12) CCHs are supported. The bit position for each CCH changes uniformly using a counter which counts from 0 through 11.
- 15 • Method 5: To support twelve (12) CCHs with pseudo-random hopping position assignments, two approaches are possible:
 - First Approach: Divide the CCHs into two groups, each containing 2^n and 2^{n-1} channels, respectively. For a 9.6 KBPS TPICCH, $n=3$. A total of five bits are derived from the long code which are referred to as long code bits 0 through 4. The position of the TCP/ISMA bit corresponding to each channel within group one varies pseudo-randomly over the first 8 bit positions in the corresponding PCG on
- 20
- 25

the TPICCH using long code bits 0 through 2. The second group is assigned the last 4 bit positions in the corresponding PCG on the TPICCH which varies pseudo randomly using long code bits 3 and 4. Long code bits 0 and 1 may be reused for the second group. However, with such reuse, undesirable dependencies may result.

- Second Approach: In order to balance the bit positioning uniformly over the total available bit positions, no grouping is performed. In this approach, 5 bits are also derived from the long code, referred to as long code bits 0 through 4. Long code bits 0 through 2 decide on a bit position within a window of 8 available positions. Long code bits 3 and 4 define an offset for starting position of the window. In this approach, the position of the TPC/ISMA bit corresponding to a given CCH is uniformly balanced over all available bit positions. Note that, it is also possible to re-use bits 0 & 1 for defining the window offset.

For a 19.2 KBPS TPICCH rate, 24 possible bit locations with a PCG are available for sending TPC and ISMA bits. According to the previously described general methodologies, the following specific implementation may be used:

- Method 1: Twenty-four (24) CCHs are supported and assigned fixed bit positions.

- Method 2: Only sixteen (16) CCHs are supported and assigned pseudo random positions using 4 bits derived from the long code mask (as in Fig. 9).
- Method 3: Twenty-four (24) CCHs are supported. The bit position corresponding to 16 CCHs are pseudo randomly changing using 4 bits derived from the long code, while the remaining 8 channels are assigned 8 fixed positions.
- Method 4: Twenty-four (24) CCHs are supported. The bit position for each CCH changes uniformly using a counter which counts from 0 through 23.
- Method 5: To support Twenty-four (24) CCHs with pseudo-random hopping position assignments, two approaches are possible:
 - First Approach: Divide the CCHs into two groups, each containing 2^n and 2^{n-1} channels, respectively. For a 19.2 KBPS TPICCH, $n=4$. A total of seven bits are derived from the long code which are referred to as long code bits 0 through 64. The position of the TCP/ISMA bit corresponding to each channel within group one varies pseudo-randomly over the first 16 bit positions in the corresponding PCG on the TPICCH using long code bits 0 through 3. The second group is assigned the last 8 bit positions in the corresponding PCG on the TPICCH which varies pseudo randomly using long

code bits 4 and 6. Long code bits 0 through 2 may be reused for the second group. However, with such reuse, undesirable dependencies may result.

- Second Approach: In order to balance the bit positioning uniformly over the total available bit positions, no grouping is performed. In this approach, 7 bits are also derived from the long code, referred to as long code bits 0 through 6. Long code bits 0 through 3 decide on a bit position within a window of 16 available positions. Long code bits 4 through 6 define an offset for starting position of the window. In this approach, the position of the TPC/ISMA bit corresponding to a given CCH is uniformly balanced over all available bit positions. Note that, it is also possible to re-use bits 0 through 2 for defining the window offset.

In General, for a CCH rate of R KBPS there are N available bit positions. It can be shown that $N = M_1 + M_2$, where $M_1 = 2^n$, $M_2 = 2^{n-1}$ and $2^{n+1} > N$. At 9.6 KBPS, $N=12$, $M_1 = 8$, $M_2 = 4$ and $n = 3$. At 19.2 KBPS, $N=24$, $M_1 = 16$, $M_2 = 8$ and $n = 4$. According to the previously described general methodologies, the following general implementation is employed:

- Method 1: N CCHs are supported and assigned fixed bit positions.

- Method 2: Only M_1 CCHs are supported and assigned pseudo random positions using n bits derived from the long code mask.
- Method 3: N CCHs are supported. The bit position corresponding to M_1 CCHs are pseudo randomly changing using n bits derived from the long code, while the remaining M_2 channels are assigned M_2 fixed positions.
- Method 4: N CCHs are supported. The bit position for each CCH changes uniformly using a counter which counts from 0 through $N-1$.
- Method 5: To N CCHs with pseudo-random hopping position assignments, two approaches are possible:
 - First Approach: Divide the CCHs into two groups, each containing 2^n and 2^{n-1} channels, respectively. A total of $2n-1$ bits are derived from the long code which are referred to as long code bits 0 through $2n-2$. The position of the TCP/ISMA bit corresponding to each channel within group one varies pseudo-randomly over the first M_1 bit positions in the corresponding PCG on the TPICCH using long code bits 0 through $n-1$. The second group is assigned the last M_2 bit positions in the corresponding PCG on the TPICCH which varies pseudo randomly using long code bits n and $2n-2$. Long code bits 0 through $2n-2$ may be reused for the second group. However, with such reuse,

undesirable dependencies may result. This particular methodology is illustrated in FIG. 13 with the power control group 1400 including group I bits 1402 and group II bits 1404.

- Second Approach: In this approach $2n-1$ bits are also derived from the long code, referred to as long code bits 0 through $2n-2$. Long code bits 0 through $2n-1$ decide on a bit position within a window of M_1 available positions. Long code bits n through $2n-2$ define an offset for starting position of the window. Note that, it is also possible to re-use bits 0 through $n-2$ for defining the window offset. This particular methodology is illustrated in FIG. 14 with the power control group 1500 shown to be randomly placed.

Of course, the examples provided above are for illustration only. Other techniques may be employed for incorporating the power control and inhibit bits into the

FIG. 11 is a block diagram illustrating the manner in which a base station may couple a quick paging channel (QPCH) and an inhibit sense multiple access channel to another forward link channel using IQ mapping. As shown, for example, the TPICCH 1102 and the QPCH may be combined onto a single Walsh channel by mapping the TPICCH to the I (in-phase) and QPCH to the Q (quadrature) portion of a Walsh channel using an IQ map 1106. Of course, the converse

mapping of these channels is also possible. Alternatively, the TPICCH channel may be mapped to I and Q or two TPICCH channels may be mapped to I and Q.

With a 19.2 KBPS rate on the TPICCH and QPCH this would require a Walsh code length of 64. The Walsh code length, however, can be chosen to match the rate of the TPICCH and QPCH to the desired chip rate. If the TPICCH and QPCH have different data rates, symbol repetition can be employed to match the rates of the channels before Walsh coding.

Furthermore, if the TPICCH channel is not transmitted, then an additional QPCH channel can be transmitted on the I (in-phase) part of the Walsh channel. In principle, a given Walsh channel may carry one QPCH, a TPICCH and QPCH, two QPCH's, or two TPICCH's. A single QPCH can be mapped to I and Q (i.e., QPSK) and a Walsh code length 2 times the length with respect BPSK (i.e., mapping to only the I or Q channel) can be accommodated.

FIG. 12 is a block diagram illustrating generally the components of a base station 1300 and of a subscriber unit 1350, each of which operates according to the present invention. While FIG. 12 shows the components of the base station 1300 and the components of the subscriber unit 1350 to be substantially similar, the construction of these devices in actuality is quite different. However, each of these devices includes the disclosed components or some variation thereof. These basic components allow the base

station 1300 and the subscriber unit 1350 to operate according to the previous description made with reference to FIGs. 1 through 12.

The base station 1300 includes processing circuitry 1302 that couples to memory 1304, that couples to a BSC interface 1306 and that couples to a coder/decoder 1308. The processing circuitry 1302 executes instructions that are stored in memory 1304 and interfaces with a coupled BSC via the BSC interface 1306. Further, the processing circuitry 1302 sends communications intended for the subscriber unit 1350 to the coder/decoder 1308.

The processing circuitry 1302 controls each of the components contained in the base station 1300 even though explicit connections are not shown from the processing circuitry 1302 to each of the other components of the base station 1300. For example, the processing circuitry 1302 directs the coder/decoder 1308 to code outgoing transmissions with a Walsh code and directs the coder/decoder 1308 to decode incoming transmissions with a particular Walsh codes. While the description of the components of these components of the base station 1300 are substantially simplified, such coding and encoding is performed on all outgoing and incoming communications, respectively.

The spreader/despreader 1310 couples to the coder/decoder 1308 and also to the Radio Frequency (RF) interface 1312. The spreader/despreader 1312 spreads

outgoing transmissions with a PN spreading code and despreads incoming transmissions with a PN spreading code. The RF interface 1312 modulates outgoing transmissions onto an RF carrier and demodulates incoming transmissions. An antenna 5 1314 couples to the RF interface and services the cell/sector(s) supported by the base station 1300.

The subscriber unit 1350 includes processing circuitry, memory 1354, a user interface and a coder/decoder 1358. Those elements in the subscriber unit 1350 having names 10 common to those elements in the base station 1300 perform similar functions. The user interface 1356 allows a user to operate the subscriber unit 1350. The coder/decoder 1358 couples to the processing circuitry 1352 as well as the spreader/despreader 1360. The spreader/despreader 1360 15 couples to an RF interface 1362, which couples to an antenna 1314.

The components of the subscriber unit 1350 may be contained in a stand-alone device. However, these components could also be contained upon a PCMCIA card that 20 couples to a laptop computer. Thus, the components of the subscriber unit 1350 illustrated support those operations previously described herein. However, these components may support additional operations as well that are not described herein.

25 FIGs. 14 and 15 illustrate the positioning of power control and inhibit bits within power control groups. The

content of these FIGs. was discussed with reference to FIG. 10. In general, the power control group positioning illustrated in FIG. 13 corresponds to method 5a, first approach of the methods described with reference to FIG. 10 while the power control group positioning in FIG. 14 corresponds to method 5b, second approach of the methods described.

While the invention is susceptible to various modifications and alternative forms, specific embodiments thereof have been shown by way of example in the drawings and detailed description. It should be understood, however, that the drawings and detailed description thereto are not intended to limit the invention to the particular form disclosed, but on the contrary, the invention is to cover all modifications, equivalents and alternatives falling within the spirit and scope of the present invention as defined by the claims.

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